L

3). A somple example & a naive approach. $f(\alpha_1) \sim N(M_1, 1)$ $\Gamma(a_1) \sim N(M_2, 1)$ costs me \$1. to pull the slot machine once ĨŁ have \$100, what is the best strategy Ι to marinize the total reward without knowly M1 8 N2 ? - Naive approach: - Try Q1 25 times - Try Q2 25 fimes - Pull the one with higher empirical means for the - rest of co times. - Why it is not optimal? - when $M_1 = 10$, $M_2 = -10$. - when $u_1 = 1$, $u_2 = -0.7$ trade off btw exploration & exploitation. - A better Algorithm: Mi is in two interval with $-VCB: UCB; = \frac{S_i}{q_i} + \sqrt{\frac{2\log(t)}{q_i}}$ t 1-5 prob ار آناز <u>s;</u> <u>q;</u> - why it is better? I when M. - My (is large when My - M2) is small.

- Wheek theoretical generates do we have?
-
$$iA = \{a_{1}, \dots, a_{n}\}$$
.
- reword of $a_{i} : X_{a_{i}} \sim \rho(X | \theta_{i})$
- Horizon : total # of pulls.
- policy: a nopping time the horizony dota \rightarrow dotribution is action space.
- policy: a nopping time the horizony dota \rightarrow dotribution is action space.
- $k_{i-1} \rightarrow A_{i}$.
- $k_{i} \rightarrow A_{i} \rightarrow A_{i} \rightarrow A_{i}$.
- $k_{i} \rightarrow A_{i} \rightarrow A_{i}$.
- $k_{i} \rightarrow A_{i} \rightarrow A_{i} \rightarrow A_{i}$.
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- $k_{i} \rightarrow A_{i} \rightarrow A_{i} \rightarrow A_{i} \rightarrow A_{i} \rightarrow A_{i}$.
- $k_{i} \rightarrow A_{i} \rightarrow$

$$\mathcal{B} \mathbb{E}[T_{i}(n)] = \mathbb{E}[T_{i}(n) | G_{i}(n)] + \mathbb{E}[T_{i}(n) | G_{i}^{i}(n)]$$

$$\leq u_{i} + \mathbb{P}(G_{i}^{i}(n)) - \frac{u_{i}c^{2}d^{2}}{2})$$

$$\leq u_{i} + \left(nd + e^{-\frac{u_{i}c^{2}d^{2}}{2}}\right)$$

$$plug m u_{i} = \left(\frac{2\log(1/d)}{(1-c)^{2}d^{2}}\right) - (if u_{i} \geq n, \quad \text{ always holds }), \quad dd = \frac{1}{n^{2}}$$

$$= \left(\frac{2\log(1/d)}{(1-c)^{2}d^{2}}\right) + if \eta^{1-2c^{2}/(1-c)^{2}}$$

$$plug m c = \frac{1}{2}$$

$$\Rightarrow \mathbb{E}[T_{i}(n)] \leq 3 + \frac{16\log(n)}{d^{2}}$$

$$\mathbb{P}_{i} = \mathbb{E}[T_{i}(n)] = \frac{3}{2}d_{i} = \frac{3}{2}d_{i} + \frac{16\log(n)}{d_{i}}$$

$$pf = (1) : Constraction: If it out is pulled u; times, then U(B; < M), it won't be publed cumy more
To T; (n) > U(i) = T; (t-1) = U; , A = i

=) U(B; (t-1, d) = U(t-1) + [2log(2)/T; (t-1) - U(t-1) + U(B)/T; (t-1) - U(t-1) + U(B)/T; (t-1) + U(B)/T; (t-1) + U(B)/T; (t-1) + U(B)/T; (t-1) + U(B)/T; (t-1), d) = A_t = curgamong U(B; (t-1, d) + i) = (X)$$

$$\begin{array}{l} \left(\begin{array}{c} \left(\begin{array}{c} \left(\begin{array}{c} \left(\begin{array}{c} \left(\right) \right) \right) \\ \left(\begin{array}{c} \left(\left(\right) \right) \\ \left(\right) \right) \\ \left(\left(\right) \right) \\ \left(\right) \\ \left($$

$$B = P\left(\hat{M}_{ini}^{*} + \sum_{i=1}^{n} 2u_{i}\right)$$

$$= P\left(\hat{M}_{ini}^{*} - M_{i}^{*} \ge M_{i} - M_{i}^{*} - \int_{u_{i}}^{1} \frac{1}{2u_{i}}u_{i}}{u_{i}}\right)$$

$$= P\left(\hat{M}_{ini}^{*} - M_{i}^{*} \ge M_{i} - M_{i}^{*} - \int_{u_{i}}^{1} \frac{1}{2u_{i}}u_{i}}{u_{i}}\right)$$

$$= P\left(\hat{M}_{ini}^{*} - M_{i}^{*} \ge M_{i} - M_{i}^{*} - \int_{u_{i}}^{1} \frac{1}{2u_{i}}u_{i}}{u_{i}}\right)$$

$$= P\left(\hat{M}_{ini}^{*} - M_{i}^{*} \ge Cd_{i}\right)$$

$$= e^{-\frac{u_{i}cd_{i}}{2}}$$

$$= e^{-\frac{u_{i}cd_{i}}{2}}$$

$$= \frac{u_{i}cd_{i}}{2}$$

Formal def
G = subguassian:
$$E = E = e^{X} = e^{\frac{X'A'}{2}}$$
 for $\forall X \in R$
 $P = P(X \ge E) \le e^{-\frac{S'}{10^{4}}}$, $P(X \le -\varepsilon) \le e^{\frac{K'}{10^{4}}}$
 $P = P(X \ge E) \le e^{-\frac{S'}{10^{4}}}$, $P(X \le -\varepsilon) \le e^{\frac{K'}{10^{4}}}$
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Time k-arned 2-autogenesisen bandet prob.
for
$$\forall$$
 horizon n, if $\delta = \frac{1}{n}$, then
 $R_n = g_{\lambda} [ritigito) + 3 \notin \Delta;$
 $Pf: R_n = \underbrace{\sum} \Delta; E[T; (n)]$
 $= \underbrace{\sum} \Delta; E[T; (n)]$
 $= \frac{\sum} \Delta; E[T; (n)]$
 $= n\Delta + \frac{k! b(g_{\Sigma}(m))}{\Delta} + \underbrace{\sum} 3b; + \frac{(bh_1(n))}{\Delta;}$
 $= n\Delta + \frac{k! b(g_{\Sigma}(m))}{\Delta} + \underbrace{\sum} 3b;$
 $\{g_{\lambda}[knlog^{(m)}] + 3 \notin \Delta;$
 $\{g_{\lambda}[knlog^{(m)}] + 3 \notin \Delta;$
 $F(\Delta n + m - c) \leq e^{-\frac{12}{kot}} = \delta = \frac{nct}{2c^k} = [og(\frac{1}{k})]$
 $=) C = \left[\frac{2C^2(b_{\Sigma}(\frac{1}{k}))}{n} + \frac{2c_{0}(\frac{1}{k})}{\pi(c_{0})} + \frac{2c_{0}(\frac{1}{k})}{\pi(c_{0})} + \frac{2c_{0}(\frac{1}{k})}{\pi(c_{0})} + \frac{2c_{0}(\frac{1}{k})}{2c_{0}} + \frac{2c_{0}(\frac{1}{k})}{2c_{0}}$